The question of Collins on the word and conjugacy problems.

Arman Darbinyan CNRS, ENS Paris

arman.darbinyan@ens.fr

GAGTA 2019 Bar Ilan University, Israel

May 30, 2019



The word and conjugacy problems (as formal languages)

Given $G = \langle S \rangle$, S is finite and symmetric (i.e. $S = S^{-1}$),

$$WP(S) = \{ w \in S^* \mid w =_G 1 \} \subseteq S^*$$

$$\operatorname{CP}(S) = \{(u, v) \in S^* \times S^* \mid u \sim_{conj} v \text{ in } G\} \subseteq S^* \times S^*$$

The word and conjugacy problems

The WP and CP for groups are important decision problems, in particular, because they serve as a bridge that connects theories of computability and computational complexity with geometric/combinatorial group theory.

The word and conjugacy problems

The WP and CP for groups are important decision problems, in particular, because they serve as a bridge that connects theories of computability and computational complexity with geometric/combinatorial group theory. For example:

Britton-Higman, 1973

A f.g. group G has a decidable word problem iff it embeds into a simple subgroup of a finitely presented group. (Later, Thompson showed that the simple subgroup can be requested to be finitely generated.)

Later, Sacerdote obtained a similar characterization for the decidability of the conjugacy problem.

Birget-Ol'shanskii-Rips-Sapir, 2001

The word problem of a f.g. group is decidable in NP time iff it embeds into a finitely presented group with polynomially bounded Dehn function.

(Un)decidability of the word and conjugacy problems

The Novikov-Boone theorem, 1955, 1958

There exists a finitely presented group with undecidable word problem.

The Novikov-Boone theorem, 1955, 1958

There exists a finitely presented group with undecidable word problem.

Miller III, 1971 and Collins, 1972

There exist groups with decidable word problem but undecidable conjugacy problem.

The Novikov-Boone theorem, 1955, 1958

There exists a finitely presented group with undecidable word problem.

Miller III, 1971 and Collins, 1972

There exist groups with decidable word problem but undecidable conjugacy problem.

The example of Miller III was finitely generated. In contrast, the example of Collins was finitely presented.

Clapham, 1967

Every f.g. group with decidable WP embeds into a finitely presented group with decidable WP.

Clapham, 1967

Every f.g. group with decidable WP embeds into a finitely presented group with decidable WP.

Olshanskii-Sapir, 2003

Every f.g. group with decidable CP embeds into a finitely presented group with decidable CP.

Collins, 1970's

Can every torsion-free f.g. group with decidable WP be embedded into a f.g. group with decidable CP?

Collins, 1970's

Can every torsion-free f.g. group with decidable WP be embedded into a f.g. group with decidable CP?

*For the case of groups with torsions a negative answer was obtained by Macintyre. The main specifics of non torsion-free case is that if two elements of a group are conjugate, then they have the same order.

Collins, 1970's

Can every torsion-free f.g. group with decidable WP be embedded into a f.g. group with decidable CP?

*For the case of groups with torsions a negative answer was obtained by Macintyre. The main specifics of non torsion-free case is that if two elements of a group are conjugate, then they have the same order. For the torsion free f.g. groups we have the following.

Osin, 2000's

Every f.g. torsion-free group embeds into a f.g. group with exactly two conjugacy classes.

For the groups with decidable power problem, Olshanskii and Sapir obtained the following positive answer to the question of Collins.

Olshanskii-Sapir, 2004

Every countable group with decidable power problem is embeddable into a 2-generated finitely presented group with decidable conjugacy and power problems. However, in general, the answer to Collins' question is negative:

D., 2017

There exists a 2-generated torsion-free group \mathcal{G} with decidable word problem that does not embed into a group with decidable conjugacy problem. Moreover, the group \mathcal{G} can be chosen to be solvable of solvability length 4 or be finitely presented. However, in general, the answer to Collins' question is negative:

D., 2017

There exists a 2-generated torsion-free group \mathcal{G} with decidable word problem that does not embed into a group with decidable conjugacy problem. Moreover, the group \mathcal{G} can be chosen to be solvable of solvability length 4 or be finitely presented.

In connection with Osin's theorem we obtain the following

Corollary

There exists a 2-generated group G with decidable word problem that does not embed into a recursively presented group with finitely many conjugacy classes.

Computable presentations for groups, Rabin, 1960, Mal'cev, 1961 The presentation $G = \langle x_1, x_2, \dots | r_1, r_2, \dots \rangle$ is called computable if for any word w from $\{x_1^{\pm 1}, x_2^{\pm 1}, \dots\}^*$, the lexicographically smallest word from the same alphabet that is equal to w in G can be computably found. Computable presentations for groups, Rabin, 1960, Mal'cev, 1961 The presentation $G = \langle x_1, x_2, \dots | r_1, r_2, \dots \rangle$ is called computable if for any word w from $\{x_1^{\pm 1}, x_2^{\pm 1}, \dots\}^*$, the lexicographically smallest word from the same alphabet that is equal to w in G can be computably found.

Example

The presentation

$$G_1 = \langle x_1, x_2, \dots \mid x_i^{-1} x_j^{-1} x_i x_j = 1, \forall i, j \in \mathbb{N} \rangle$$

is a computable presentation for the free abelian group of countable rank $\bigoplus \mathbb{Z}$. Now let us fix a non-recursive set $\mathcal{N} \subset 2\mathbb{N}$ and consider the following presentation

$$G_2 = \langle x_1, x_2, \dots \mid [x_i, x_j] = 1, \forall i, j \in \mathbb{N} \& x_{2i} = x_{2i+1} \text{ iff } i \in \mathcal{N} \rangle.$$

The last presentation is not computable but still it is a presentation of the group $\bigoplus \mathbb{Z}$.

Embedding theorem for groups, D., 2015

Let $G = \langle X \rangle$ be a group with countable generating set $X = \{x_1, x_2, \ldots\}$. Then there exists an embedding $\Phi_X : G \hookrightarrow K$ into a two-generated group $K = \langle f, c \rangle$ such that:

- K has a recursive presentation if and only if G has a recursive presentation with respect to the generating set X;
- 2 K has decidable word problem if and only if G is computable with respect to the generating set X;
- 3 If $X = \{x_1, x_2, \ldots\}$ is recursively enumerated, then there exists a computable map $\phi_X : i \mapsto \{f^{\pm 1}, s^{\pm 1}\}^*$ such that ϕ_X represents the element $\Phi_X(x_i)$ in K;
- 4 There exists $N \lhd K$ such that $\Phi_X(G) \lhd N$, and K/N, $N/\Phi_X(G)$ are abelian groups;
- **5** The membership problem for the subgroup $\Phi_X(G) \leq K$ is decidable.





{n_1, n_2, ...} and {m_1, m_2, ...} are recursively enumerable and recursively inseparable subsets of N

Thank you!

◆□▶ ◆□▶ ◆臣▶ ◆臣▶ 臣 のへで