### Trusses

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#### References:

 TB, Trusses: between braces and rings, arXiv:1710.02870 (2017)

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► TB, *Towards semi-trusses*, Rev. Roumaine Math. Pures Appl. (Tome LXIII No. 2, 2018)

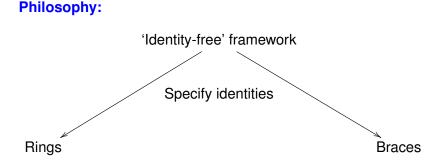
# Aim and philosophy:

**Aim:** To present an algebraic framework for studying braces and rings on equal footing.



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Herds (or heaps or torsors)

H. Prüfer (1924), R. Baer (1929)

#### Definition

A *herd* (or *heap* or *torsor*) is a nonempty set A together with a ternary operation

$$[-, -, -]$$
:  $A \times A \times A \rightarrow A$ ,

such that for all  $a_i \in A$ ,  $i = 1, \ldots, 5$ ,

$$[[a_1, a_2, a_3], a_4, a_5] = [a_1, a_2, [a_3, a_4, a_5]],$$

$$[a_1, a_2, a_2] = a_1 = [a_2, a_2, a_1].$$

A herd (A, [-, -, -]) is said to be *abelian* if

 $[a, b, c] = [c, b, a], \text{ for all } a, b, c \in A.$ 

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# Herds are in '1-1' correspondence with groups

If (A, ◊) is a (abelian) group, then A is a (abelian) herd with operation

$$[a, b, c]_\diamond = a \diamond b^\diamond \diamond c.$$

▶ Let (A, [-, -, -]) be a (abelian) herd. For all  $e \in A$ ,

$$a\diamond_e b:=[a,e,b],$$

makes *A* into (abelian) group (with identity *e* and the inverse mapping  $a \mapsto [e, a, e]$ .)

- ► Note:
  - different choices of *e* yield different albeit isomorphic groups.

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irrespective of e: [a, b, c]<sub>◊e</sub> = [a, b, c].

# Herds are 'groups without specified identity'

There if a forgetful functor

 $\text{Grp} \longrightarrow \text{Set}_*.$ 

Morphisms from (A, [−, −, −]) to (B, [−, −, −]) are functions f : A → B respecting ternary operations:

f([a, b, c]) = [f(a), f(b), f(c)].

There is a forgetful functor

 $Hrd \longrightarrow Set$ ,

but not to the category of based sets.

• Worth noting:

$$\operatorname{Aut}(A, [-, -, -]_\diamond) = \operatorname{Hol}(A, \diamond).$$

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Worth noting:

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### Trusses

A left skew truss is a herd (A, [−, −, −]) together with an associative operation · that left distributes over [−, −, −], i.e.,

$$a \cdot [b, c, d] = [a \cdot b, a \cdot c, a \cdot d].$$

- If (A, [-, -, -]) is abelian, then we have a *left truss*.
- Right (skew) trusses are defined similarly.
- ► A truss is a triple (A, [-, -, -], ·) that is both left and right truss.

A morphism of (left/right skew) trusses is a function preserving both the ternary and binary operations.

### Trusses: between braces and (near-)rings

#### Let $(A, [-, -, -], \cdot)$ be a left skew truss.

► Assume that (A, ·) is a group with a neutral element e. Then (A, ◇<sub>e</sub>, ·) is a left skew brace, i.e.

$$a \cdot (b \diamond_e c) = (a \cdot b) \diamond_e a^{\diamond_e} \diamond_e (a \cdot c).$$

• Assume that  $e \in A$  is such that

 $a \cdot e = e$ , for all  $a \in A$ .

Then  $(A, \diamond_e, \cdot)$  is a left near-ring, i.e.

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### Trusses: generalised distributivity

Let  $(A, \diamond)$  be a group and  $(A, \cdot)$  be a semigroup. TFAE:

• There exists  $\sigma : A \rightarrow A$ , such that

$$a \cdot (b \diamond c) = (a \cdot b) \diamond \sigma(a)^{\diamond} \diamond (a \cdot c).$$

• There exists  $\lambda : A \times A \rightarrow A$ , such that,

$$a \cdot (b \diamond c) = (a \cdot b) \diamond \lambda(a, c).$$

• There exists  $\mu : A \times A \rightarrow A$ , such that

$$a \cdot (b \diamond c) = \mu(a, b) \diamond (a \cdot c)$$

• There exist  $\kappa, \hat{\kappa} : A \times A \rightarrow A$ , such that

$$a \cdot (b \diamond c) = \kappa(a, b) \diamond \hat{\kappa}(a, c).$$

•  $(A, [-, -, -]_\diamond, \cdot)$  is a left skew truss.

## Trusses from split-exact sequences of groups

Let (A, ◊) be a middle term of a split-exact sequence of groups

$$1 \longrightarrow G \longrightarrow A \xrightarrow{\alpha}_{\beta} H \longrightarrow 1$$

Let · be an operation on A defined as

$$a \cdot b = a \diamond \beta(\alpha(b))$$
 or  $a \cdot b = \beta(\alpha(a)) \diamond b$ .

• Then  $(A, [-, -, -]_{\diamond}, \cdot)$  is a left skew truss.

### The endomorphism truss

- Let (A, [-, -, -]) be an abelian herd.
- Set  $\mathcal{E}(A) := \text{End}(A, [-, -, -]).$
- $\mathcal{E}(A)$  is an abelian herd with inherited operation

[f, g, h](a) = [f(a), g(a), h(a)].

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▶  $\mathcal{E}(A)$  together with [-, -, -] and composition  $\circ$  is a truss.

### Notes on the endomorphism truss:

- Choosing the group structure f ◇<sub>id</sub> g on E(A), we obtain a two-sided brace-type distributive law between ◇<sub>id</sub> and ○.
- Fix e ∈ A, and let ε : A → A, be given by ε : a → e. Then ε ∈ ε(A), and choosing the group structure f ◊<sub>ε</sub> g on ε(A) we get a ring (ε(A), ◊<sub>ε</sub>, ◦).
- The left multiplication map

$$\ell: A \to \mathcal{E}(A), \qquad a \mapsto [b \mapsto a \cdot b],$$

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is a morphism of trusses.

Many technics and constructions familiar in ring theory can be applied to trusses.

• An *ideal* of  $(A, [-, -, -], \cdot)$  is a sub-herd X such that,

 $a \cdot x, x \cdot a \in X$ , for all  $x \in X, a \in A$ .

▶ *X* defines an equivalence relation, for  $a, b \in A$ ,

 $a \sim_X b$  iff  $\exists x \in X, [a, b, x] \in X$ .

• The quotient  $A/X := A/\sim_X$  is a truss with operations

 $[\overline{a}, \overline{b}, \overline{c}] = \overline{[a, b, c]}, \quad \overline{a} \cdot \overline{b} = \overline{a \cdot b}.$ 

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# Products, functions, polynomials

The product of (skew) trusses and the mapping truss Map(X, A) can be defined in ways analogous to that for rings, e.g. for all f, g, h ∈ Map(X, A), x ∈ X,

 $[f, g, h](x) = [f(x), g(x), h(x)], \quad (f \cdot g)(x) = f(x) \cdot g(x).$ 

Given a (commutative) truss A, a formal series truss A[[x]] is the function truss

$$A[[x]] := \operatorname{Map}(\mathbb{N}, A).$$

► For an idempotent element e of (A, ·) one can define e-polynomial truss A<sub>e</sub>[x] by

 $A_{e}[x] := \{f \in A[[x]] \mid f(i) \neq e \text{ for finitely many } i \in \mathbb{N}\}.$ 

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### Modules of trusses

► A left module over a truss (A, [-, -, -], ·) is an abelian herd (M, [-, -, -]) together with a morphism of trusses

$$\pi_{\boldsymbol{M}}: \boldsymbol{A} \to \mathcal{E}(\boldsymbol{M}).$$

The action of A on M, a ▷m := π<sub>M</sub>(a)(m), satisfies: Distributive laws:

$$a \triangleright [m_1, m_2, m_3] = [a \triangleright m_1, a \triangleright m_2, a \triangleright m_3],$$
$$[a, b, c] \triangleright m = [a \triangleright m, b \triangleright m, c \triangleright m],$$

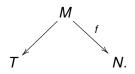
Associative law:

$$a \triangleright (b \triangleright m) = (a \cdot b) \triangleright m.$$

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# Category of modules

- Morphisms of modules over trusses are defined as functions preserving the ternary operations and actions; category A – Mod.
- Right modules, bimodules defined analogously.
- ► A Mod has a terminal object T = {0} but not an initial object.
- A Mod has cokernels, i.e. pushouts of



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# Category of modules

- A Mod has quotients:
  - ► Take a submodule *N* of *M*.
  - Define an equivalence relation, for  $m_1, m_2 \in M$ ,

 $m_1 \sim_N m_2$  iff  $\exists n \in N, [m_1, m_2, n] \in N$ .

• 
$$\overline{M} := M/N := M/\sim_N$$
,

 $[\overline{m_1},\overline{m_2},\overline{m_3}]=\overline{[m_1,m_2,m_3]}, \quad a\triangleright\overline{m}=\overline{a\triangleright m}.$ 

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• Given a morphism of A-modules  $f: M \rightarrow N$ ,

$$\operatorname{coker}(f) = N/\operatorname{Im}(f).$$

## Endomorphism and matrix trusses

► For any *A*-module *M*,

#### $\operatorname{End}_A(M)$

is a truss in the same way as endomorphisms of an abelian herd.

A<sup>n</sup> is an A-module: for all a = (a<sub>i</sub>), b = (b<sub>i</sub>), c = (c<sub>i</sub>) ∈ A<sup>n</sup>, x ∈ A,

$$[a,b,c]_i = [a_i,b_i,c_i], \qquad (x \triangleright a)_i = x \triangleright a_i.$$

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- $M_n(A) := \operatorname{End}_A(A^n)$  is a (matrix) truss.
- End<sub>A</sub>(A<sup>n</sup>) satisfy a brace-type distributive law between ⊲<sub>id</sub> and ∘.

Although this is the end of the story so far it might be just beginning...

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