# PRIMITIVE ALGEBRAS WITH ARBITRARY GELFAND-KIRILLOV DIMENSION

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ABSTRACT. We construct, for every real  $\beta \geq 2$ , a primitive affine algebra with Gelfand-Kirillov dimension  $\beta$ . Unlike earlier constructions, there are no assumptions on the base field. In particular, this is the first construction over  $\mathbb{R}$  or  $\mathbb{C}$ .

Given a recursive sequence  $\{v_n\}$  of elements in a free monoid, we investigate the quotient of the free associative algebra by the ideal generated by all non-subwords in  $\{v_n\}$ .

We bound the dimension of the resulting algebra in terms of the growth of  $\{v_n\}$ . In particular, if  $|v_n|$  is less than doubly-exponential then the dimension is 2. This also answers affirmatively a conjecture of Salwa [12].

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# 1. Preliminaries

Let A be an affine k-algebra. The **Gelfand-Kirillov dimension** [6] of A is defined as

$$GKdim(A) = \limsup_{s \to \infty} \frac{\log \dim(V + V^2 + V^3 + \dots + V^s)}{\log s}$$

where V is a finite-dimensional subspace that generates A as an algebra. (see [9] for details).

It is easily seen that GKdim(A) = 0 iff A is finite dimensional. Otherwise  $GKdim(A) \ge 1$ , and by Bergman's gap theorem [3], either GKdim(A) = 1 (in which case A is a PI-algebra by [14]), or  $GKdim(A) \ge 2$ .

If A is PI then GKdim(A) equals the transcendence degree of A over k [2], and is thus an integer.

Affine algebras with GKdim arbitrary real  $\beta \geq 2$  were constructed by Borho-Kraft [4] and by Warfield [15] (*cf.* also [9, 2.9]). These examples fail to be prime. For a semiprime example, see [7].

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In [8] Irving and Warfield constructed primitive algebras with arbitrary GKdim, under the restriction that the base field F has an infinite-dimensional algebraic extension.

In this note we provide straightforward examples of primitive affine algebras over an arbitrary field, having arbitrary  $GKdim \geq 2$ .

Our construction is a generalization of the Morse algebra. It is a monomial algebra, that is, the quotient of a free associative algebra by an ideal generated by monomials. The growth of such algebras has been studied, for example, in [1], [5] and [11].

We assume the ideal to be generated by all monomials that are not subwords in a given sequence  $\{v_n\}$  of elements in a free monoid, and relate properties of the resulting algebra A to the sequence  $\{v_n\}$ .

In section 2 we prove that under conditions (1), (2) below, GKdim(A) is bounded in terms of  $\frac{\log |v_{n+1}|}{\log |v_n|}$ . If  $|v_n|$  is less then doubly-exponential (as is the case if  $v_n$  is defined by a constant recursion rule), then  $GKdim(A) \leq 2$ . Some theory of recurring sequences over finite fields is used in section 3 to choose  $\{v_n\}$  that achieve the bound, thus producing prime algebras of arbitrary GKdim. These examples are shown to be primitive in section 4.

# 2. Monomial Algebras

Let S be a free finitely-generated monoid. If L is an ideal of S, then  $A = k[S/L] \cong k[S]/k[L]$  is an affine **monomial algebra**.

Fix a group  $\mathcal{T}$  of permutation automorphisms of  $\mathcal{S}$ . If  $M < \mathcal{S}$  is a submonoid, let  $M_{\mathcal{T}}$  denote the closure of M under the operation of  $\mathcal{T}$ .

Let  $\{v_n\}$  be a sequence in S. Evidently, the set L of words that are not subwords in any  $v_n$ , is an ideal. In this case {subwords of length s in  $v_n$ : all n} is a basis for the s'th homogeneous part of A.

All algebras discussed below are defined over a fixed (but arbitrary) field k, with ideals L as above, where we assume

(1) 
$$\forall n : |v_n| \le |v_{n+1}|$$

(2) 
$$\forall n : v_{n+\kappa} \in \langle v_n, v_{n+1}, \dots, v_{n+\kappa-1} \rangle_{\mathcal{T}}$$

for some fixed  $\kappa$ .

Write  $x \leq y$  if x is a subword in y. As in Salwa [12], there is an obvious criterion for A to be prime.

**Remark 2.1.** A is prime iff for any i, j there exist  $w \in S$  such that  $v_i w v_j < v_k$  for some k.

If  $x \in \mathcal{S}$ ,  $X \subseteq \mathcal{S}$ , let  $W_s(x) = \{w : w \leq x, |w| = s\}$  and  $W_s(X) = \bigcup_{x \in X} W_s(x)$ . Obviously,  $|W_s(x)| < |x|$  if s > 1. Write  $w_s = |W_s(\{v_n\})|$ . Then

$$GKdim(A) = \limsup_{s \to \infty} \frac{\log(w_1 + w_2 + \dots + w_s)}{\log s}.$$

Note that GKdim(A) = 0 iff  $|v_n|$  is bounded. We assume henceforth that this is *not* the case, so by Bergman's gap theorem GKdim(A) = 1 or  $GKdim(A) \ge 2$ .

The main tool we use to compute GKdim(A) is the following simple lemma.

# **Lemma 2.2.** If $|x_2|, |x_3|, \ldots, |x_{m-1}| \ge s$ , then

$$W_s(x_1x_2...x_m) = W_s(x_1x_2) \cup W_s(x_2x_3) \cup \cdots \cup W_s(x_{m-1}x_m).$$

*Proof.* A subword of length s of  $x_1x_2...x_m$  can never intersect more then two consecutive  $x_i$ 's.

**Theorem 2.3.** Let  $d = \limsup \frac{\log |v_{n+1}|}{\log |v_n|}$ . Then  $GKdim(A) \le 1 + d^{\kappa}$ .

*Proof.* Fix s and some  $\epsilon > 0$ .

There is some  $\mu$  such that  $|v_{\mu}| < s \leq |v_{\mu+1}|$ . Iterating assumption (2), we get  $v_{\mu+i} \in \langle v_{\mu+1}, \ldots, v_{\mu+\kappa} \rangle_{\mathcal{T}}$  for all  $i \geq 1$ . Thus, by Lemma 2.2,

$$w_{s} = |W_{s}(\{v_{\mu+i} : i \geq 1\})|$$

$$\leq |\bigcup_{\tau_{1},\tau_{2}\in\mathcal{T}, 0 < j_{1}, j_{2} \leq \kappa} W_{s}(\tau_{1}(v_{\mu+j_{1}})\tau_{2}(v_{\mu+j_{2}}))|$$

$$\leq \sum_{\tau_{1},\tau_{2}\in\mathcal{T}, 0 < j_{1}, j_{2} \leq \kappa} |W_{s}(\tau_{1}(v_{\mu+j_{1}})\tau_{2}(v_{\mu+j_{2}}))|$$

$$< \sum_{\tau_{1},\tau_{2}\in\mathcal{T}, 0 < j_{1}, j_{2} \leq \kappa} |\tau_{1}(v_{\mu+j_{1}})\tau_{2}(v_{\mu+j_{2}})|$$

$$\leq 2|\mathcal{T}|^{2}\kappa^{2}|v_{\mu+\kappa}|$$

Let  $c = 2|\mathcal{T}|^2 \kappa^2$ . We have that  $w_1 + w_2 + \dots + w_s < cs|v_{\mu+\kappa}|$ , so  $\frac{\log(w_1 + \dots + w_s)}{\log s} < \frac{\log cs}{\log s} + \frac{\log |v_{\mu+\kappa}|}{\log |v_{\mu}|} \le 1 + d^{\kappa} + \epsilon$  for large enough s.

If  $|v_n|$  is less than doubly exponential, *i.e.*,  $\log \log |v_{n+1}| - \log \log |v_n| \rightarrow 0$ , then d = 1 and  $GKdim(A) \leq 2$ .

It can be shown that GKdim(A) = 1 iff for some constant C, almost all the words  $v_n$  are periodic with period < C. We omit the details of the proof.

In many natural examples  $v_n$  are defined recursively. In this case we have

**Corollary 2.4.** Suppose that  $\{v_n\}$  is defined by a constant recursion rule (i.e. the formula for  $v_n$  as a function of  $v_{n-1}, \ldots, v_{n-\kappa}$  does not involve n), such that assumption (1) is satisfied.

Then  $|v_{n+1}| < M|v_n|$  for some constant M, and by Theorem 2.3 we have that GKdim(A) < 2.

In particular Salwa's example [12] has Gelfand-Kirillov dimension 2. We end this section with

**Lemma 2.5.** Assume that for any  $i, v_i < v_k$  for some k. Then  $GKdim(A) \ge 1 + limsup \frac{\log w_s}{\log s}.$ 

*Proof.* The assumption implies that  $w_s$  is nondecreasing. Now

$$GKdim(A) = limsup \frac{\log(w_1 + \dots + w_{2s})}{\log 2s}$$
  

$$\geq limsup \frac{\log sw_s}{\log 2s}$$
  

$$= 1 + limsup \frac{\log w_s}{\log s}.$$

### 3. PRIME AFFINE ALGEBRAS WITH ARBITRARY DIMENSION

In this section we present sequences  $\{v_n\}$  that define prime algebras with arbitrary GKdim > 2. These examples are shown to be primitive in Section 4.

Some preliminaries from the theory of linear recurring sequences are needed. The reader is referred to [10, Chap. 8] for more details and proofs.

**Proposition 3.1.** Let m > 1 be a natural number. Let K be the field of order  $2^m$ . Pick a generator u of the multiplicative group  $K^*$ . Let  $g(x) = g_0 + g_1 x + \cdots + x^m$  be the minimal polynomial of u over  $\mathbb{Z}_2$ .

Define a sequence  $\{b_i\}$  over  $\mathbb{Z}_2$  by  $b_0 = \cdots = b_{m-2} = 0, b_{m-1} = 1$ , and the recursion rule  $b_{i+m} = g_0 b_i + g_1 b_{i+1} + \cdots + g_{m-1} b_{i+m-1}$   $(i \ge 0)$ . Then  $\{b_i\}$  has period  $2^m - 1$ , and for every non-zero  $w \in \mathbb{Z}_2^m$ , there

is a unique  $0 \le i < 2^m - 1$  such that  $w = b_i b_{i+1} \dots b_{i+m-1}$ .

Moreover, if w and w' are opposite non-zero words (i.e. w + w' =(11...1) of length m+1, then exactly one of them appears in  $\{b_i\}$ ((11...1), (00...0) do not appear at all).

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**Definition 3.2.**  $L_m$  denotes a word of length  $2^m + m - 1$  over  $\mathbb{Z}_2$ , constructed as the first  $2^m + m - 2$  elements of a sequence defined as in Proposition 3.1, preceded by a single 0.

For example, we can take  $L_1 = 01$ ,  $L_2 = 00110$ ,  $L_3 = 0001011100$ and  $L_4 = 0000100110101111000$ .

**Remark 3.3.** Every word of length m appears exactly once as a subword of  $L_m$ . Two opposite words of length m + 1 do not appear both in  $L_m$  except for the couple (00...01), (11...10). No opposite words of length m + 2 appear in  $L_m$ .

Let  $S = \langle x, y \rangle$  be the free monoid on two generators, with the automorphism  $v \mapsto \bar{v}$  defined by  $\bar{x} = y, \bar{y} = x$ . The substitution of a word v in  $L_m$  is defined as the replacement of all 0's in  $L_m$  by v and all 1's by  $\bar{v}$ . For example,  $L_1(v) = v\bar{v}$ .

Fix a sequence of integers  $r_n$ . We define  $\{v_n\} \subseteq S$  as follows:

$$v_1 = x, \quad v_{n+1} = L_{r_n}(v_n).$$

Define an algebra A using  $\{v_n\}$  as in the beginning of section 2. Note that assumptions (1) and (2) are satisfied (with  $\kappa = 1$ ).

Note that if  $r_n > 1$  then  $v_{n+1}$  starts with  $v_n v_n$  and ends with  $v_n$ . The case  $r_n = 1$  gives the well known Morse algebra.

## **Theorem 3.4.** A is prime.

*Proof.* By Remark 2.1 we must show for any  $v_i, v_j$  that  $v_i w v_j < v_k$  for some k and a word w. Pick  $n = max\{i, j\}$ , then  $v_i, v_j < v_n$ , so pick  $w_i, w_j$  such that  $v_i w_i$  is a tail of  $v_n$  and  $w_j v_j$  a head of  $v_n$ .

If  $r_n > 1$  then  $v_i w_i w_j v_j < v_n v_n < v_{n+1}$  by definition of  $L_{r_n}$ . Otherwise suppose  $r_n = 1$ ; if  $r_{n+1} > 1$  then  $v_i w_i \overline{v_n} w_j v_j < v_{n+1} v_{n+1} < v_{n+2}$ , and if  $r_{n+2} = 1$  then  $v_i w_i \overline{v_n} w_j v_j < v_{n+1} \overline{v_{n+1}} = v_{n+2}$ .

From now on we assume that  $r_n \geq 3$ .

**Lemma 3.5.** If  $m \ge (r_n + 2)|v_n|$ , then the subwords of length m in  $v_{n+1}$  and in  $\overline{v_{n+1}}$  are all different.

*Proof.* Let  $k = r_n + 2$ . It is enough to prove the assertion in the case  $m = k|v_n|$ . For n = 1 the result follows from 3.3. Let n > 1.

Let a, b be two equal subwords of  $v_{n+1}$  or  $\overline{v_{n+1}}$ . Write  $v_{n+1}$  as a word on the letters  $v_n, \overline{v_n}$  which we call *full letters*. Then a, b are determined by the full letter  $(v_n \text{ or } \overline{v_n})$  in  $v_{n+1}$  or in  $\overline{v_{n+1}}$  in which they start, and the relative position in this full letter. The strategy is to show first that a and b start at the same relative position, and then show that they actually start at the same full letter.

Write  $a = a_0 u_1 \dots u_{k-1} a_1$  where  $|a_0|, |a_1| \leq |v_n|$ , and each  $u_i$  equals one of the full letters  $v_n, \overline{v_n}$ ; write  $b = b_0 w_1 \dots w_{k-1} b_1$  in the same manner. W.l.o.g. we assume  $|a_0| \ge |b_0|$ .

Write  $a_0 = a_{00}a_{01}$  and  $b_1 = b_{10}b_{11}$  where  $a_{00} = b_0$  and  $b_{11} = a_1$ . Also factor  $u_i = u'_i u''_i$  and  $w_i = w''_i w'_i$  where  $|u''_i| = |w''_i| = |a_0| - |b_0|$ .

Assume  $|a_0| - |b_0| \leq \frac{1}{2} |v_n|$  (the other case is treated similarly). Then

 $u'_1 = w'_1$  is an equality of words of length  $\geq \frac{1}{2}|v_n|$ . Since  $\frac{1}{2}|v_n| \geq \frac{r_{n-1}+2}{2^{r_{n-1}}+r_{n-1}-1}|v_n| = (r_{n-1}+2)|v_{n-1}|$ , the induction hypothesis force  $u'_i, w'_i$  to begin in the same relative position. But then it follows that  $u''_i, w''_i$  are empty words and each of  $u'_i, w'_i$  is a full letter,  $v_n$  or  $\overline{v_n}$ . By Remark 3.3, the  $r_n + 1$  equalities  $u_i = u'_i = w'_i = w_i$  force one of two cases: a, b begin in the same position in  $v_{n+1}$  or in  $\overline{v_{n+1}}$ , in which case we are done, or  $u_1 \dots u_{k-1} = w_1 \dots w_{k-1} = v_n \dots v_n \overline{v_n}$  and  $a \leq v_{n+1}, b \leq \overline{v_{n+1}}$  (or vice versa). But  $v_n \dots v_n \overline{v_n}$  is the header of  $v_{n+1}$ , so  $a_0, b_0$  must be empty. Then we have that  $a_1 = b_1$ , the  $(r_n + 2)$ 'th equality of full letters, a contradiction of Remark 3.3.

We can now compute the Gelfand-Kirillov dimension of A.

**Theorem 3.6.** Let  $d = \limsup \frac{r_1 + \dots + r_n}{r_1 + \dots + r_{n-1}}$ . Then GKdim(A) = d + 1. *Proof.* Note that  $|v_n| = |L_{r_1}| |L_{r_2}| \dots |L_{r_{n-1}}|$ , so

$$r_1 + \dots + r_{n-1} < \log_2 |v_n| < n + r_1 + \dots + r_{n-1}.$$

By Theorem 2.3 we have

$$GKdim(A) \le 1 + \limsup \frac{\log_2 |v_{n+1}|}{\log_2 |v_n|} = 2 + \limsup \frac{\log_2 |L_{r_n}|}{\log_2 |v_n|} \le 2 + \limsup \frac{r_n + 1}{r_1 + \dots + r_{n-1}} = 1 + d.$$

For the other direction, recall that by Lemma 3.5 all of the subwords of length  $s = (r_n + 2)|v_n|$  in  $v_{n+1}$  are different. Thus

$$w_s \ge |v_{n+1}| - s = (2^{r_n} - 3)|v_n|$$

(where  $w_s$  is the number of subwords of length s in any  $v_n$ ), and, if d > 1,

$$\frac{\log_2 w_s}{\log_2 s} > \frac{\log_2(2^{r_n} - 3) + \log_2 |v_n|}{\log_2(r_n + 2) + \log_2 |v_n|} > \frac{r_1 + \dots + r_n - 1}{\log_2(r_n) + n + r_1 + \dots + r_{n-1}}.$$

The *limsup* of the lower bound is d (whether d is finite or infinite). If d = 1, then the expression in the middle already apporaches 1.  $\square$ 

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Finally, let  $\beta \in \mathbb{R}, \beta \geq 1$ .

Take  $r_n = \max([\beta^n], 3)$ , and define an algebra A as above. Checking the conditions of Theorem 3.6, we arrive at

**Theorem 3.7.** A is an affine prime algebra with  $GKdim(A) = \beta + 1$ .

In particular, the bound in 2.3 is tight (at least for  $\kappa = 1$ ).

## 4. OUR EXAMPLES ARE PRIMITIVE

In this section we show that A is primitive. We assume that  $r_n \ge 3$ , and  $r_n > 3$  infinitely often (note that for dimension 2 we must take  $r_n = 4$ ).

**Definition 4.1.** For  $u, v \in S$ , let  $u \leq_l v$  ( $u \leq_r v$ ) denote that u is a head (tail) of v. An element  $a \in A$  is a **left** (**right**) **tower** if the set of monomials of a is linearly ordered by  $\leq_l (\leq_r)$ .

Being a left tower is invariant under multiplication by a monomial from the left.

**Lemma 4.2.** Let L < A be a left ideal. Then L contains a left tower. Moreover, for every  $a \in A$ , wa is a left tower for some monomial w.

*Proof.* It is enough to show that if  $a_1, a_2$  are monomials, and  $wa_1 = 0$  iff  $wa_2 = 0$  (all  $w \in S$ ), then  $a_1, a_2$  are  $\leq_l$ -comparable. Assume  $|a_1| \geq |a_2|$ .

For some n we have  $a_2 \leq v_n$  and  $r_n > 3$ . Write  $v_n = \alpha a_2 \beta$ ,  $v_{n+1} = uv_n$ .

By assumption  $u\alpha a_1 \neq 0$ , so  $u\alpha a_1 \leq v_m$  for some m. Writing  $v_m$  as a word in the letters  $v_{n+1}, \overline{v_{n+1}}$ , the intersection of u with some letter is of length  $> \frac{1}{2}|u| > (r_n+2)|v_n|$ . By Lemma 3.5, u (and thus  $u\alpha a_1$ ) appears in  $v_m$  as a header of  $v_{n+1}$ , and we get  $u\alpha a_2 \leq_l u\alpha a_1$ , as desired.  $\Box$ 

# **Corollary 4.3.** Let $I \trianglelefteq A$ be an ideal. Then I contains a monomial.

*Proof.* By Lemma 4.2 and left-right symmetry, there is some  $a \in I$  that is a left and right tower.

Let u, w be two different monomials in  $a, |u| \leq |w|$ . Multiplying by long enough monomials from both sides we may assume that  $w = v_n$ , and  $\frac{1}{2}|v_n| < |u|$ .

Now  $u \leq_r v_n$  and  $u \leq_l v_n$  implies by Lemma 3.5 that  $u = v_n$ , a contradiction.

Let  $J = \langle x, y \rangle \trianglelefteq A$ , a maximal ideal in A.

Recall that if  $J_1, \ldots, J_n \subseteq R$  are maximal ideals in a (unital) ring R, and  $J_1 J_2 \ldots J_n \subseteq I \neq R$ , then  $I \subseteq J_i$  for some i (for taking a maximal ideal  $I_1 \supseteq I$  we have some  $J_i \subseteq I_1$ , implying  $J_i = I_1 \supseteq I$ ).

**Corollary 4.4.** J is the unique maximal ideal in A. Moreover, for any  $I \trianglelefteq A$ , the quotient A/I is a finite dimensional k-algebra and Jac(A/I) = J/I.

*Proof.* By Corollary 4.3,  $v_n \in I$  for some n. Now every monomial of length  $> 2|v_{n+1}|$  contains  $v_{n+1}$  or  $\overline{v_{n+1}}$  and thus  $v_n$  as a subword, so  $J^m \subseteq I$  for  $m = 2|v_{n+1}|$ , and thus  $I \subseteq J$ .

A/I is finite dimensional spanned by { words of length < m }, and  $Jac(A/I) = Jac(A/J^m)/(I/J^m) = (J/J^m)/(I/J^m) = J/I$ .

**Corollary 4.5.** The only prime ideals of A are 0, J. In particular, the (classical) Krull dimension of A is one.

*Proof.* 0 is prime by Remark 2.1. If  $I \neq 0$  is prime then J/I = Jac(A/I) = 0 so I = J.

Theorem 4.6. A is primitive.

*Proof.* Assume, on the contrary, that A is non-primitive. Then the only primitive ideal of A is J, so this is the Jacobson radical of A. In particular,  $x + y \in J$  should be quasi-regular.

But if a(1-x-y) = 1, let w be a longest monomial in a. Necessarily  $wx \neq 0$  or  $wy \neq 0$ , so wx (say) appears on the left hand side of the equality but not on the right hand side, a contradiction.

From Corollary 4.4 it follows that A satisfies ACC on two-sided ideals (since a chain going up from  $I \trianglelefteq A$  has length  $< \dim(A/I)$  as a k-algebra). For one-sided ideals the situation is not that pleasant.

# **Remark 4.7.** A is a not left-Noetherian.

*Proof.* We construct a left ideal that is not finitely-generated.

Ordering the monomials in A by  $\langle r \rangle$  we get a tree  $T_1$ , in which every node is a root for a tree with at least two, and thus infinitely many, branches (indeed, if  $w \langle v_n \rangle$  is a monomial write  $v_n = \alpha w \beta$ , then  $w \langle r \rangle v_n \alpha w$  and  $w \langle r \rangle \overline{v_n} \alpha w$ ). Now define a tree  $T_n$  and  $w_n \in T_n$  $(n \geq 1)$  as follows: Let  $w \in T_n$  be a monomial with  $wx, wy \in T_n$ . Take  $w_n = wx$ , and  $T_{n+1}$  the tree with root wy.

Then  $L = Aw_1 + Aw_2 + \ldots$  is definitely not *f.g.*, since the  $w_i$  are  $\leq_l$ -incomparable.

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